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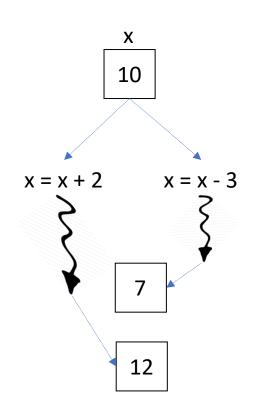
OpenMP Advanced Concepts

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OpenMP – Race Conditions

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- Inadequate management of concurrent accesses may create wrong results
 - Data races
 - Read-after-write and write-after-read dependencies must be handled by the programmer
- OpenMP provides pragmas to control access to shared data
 #pragma omp critical
- Algorithms and/or data structures may need to be redesigned

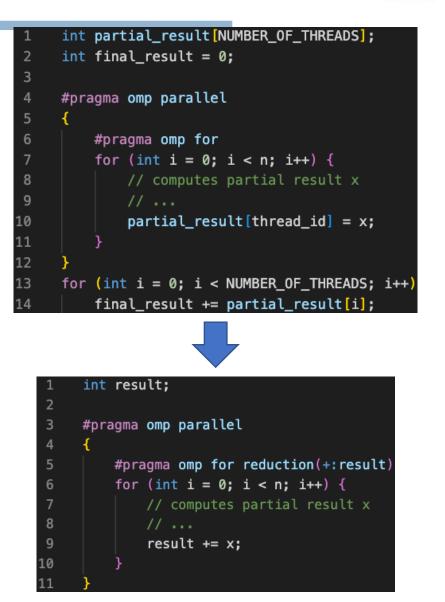


OpenMP – Reduction

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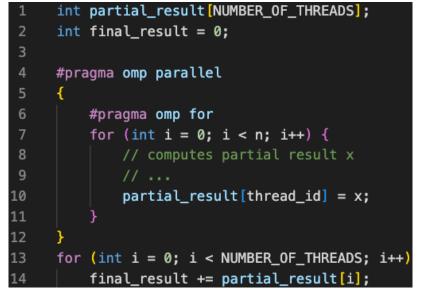
Threads often compute a subset of a workload

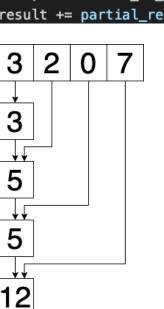
- Partial results stored privately to each thread
 - Improves performance (less contention to access shared data)
 - Reduces the risk of data races
- Partial results must be merged after the parallel section
 - can be implemented manually
 - #pragma omp reduction does this automatically (restrictions apply)

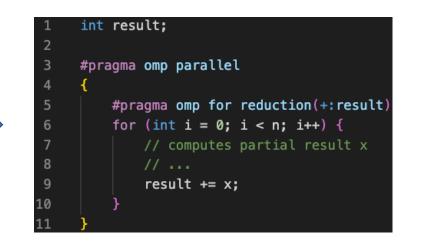


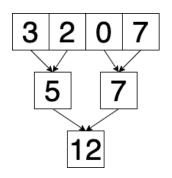
OpenMP – Reduction

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Controlling the sections of code that each thread executes can be achieved through:

- Each thread individual id
- OpenMP pragma directives
 - Master: only the master thread executes the code section
 - Single: the first thread to arrive executes the code section

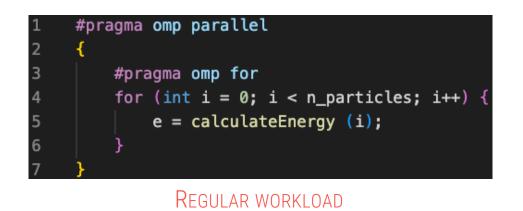


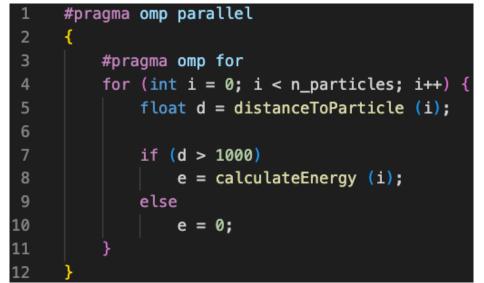
OpenMP – Loop Scheduling

Workloads can be classified as

- Regular: each loop iteration takes the same amount of time
- Irregular: loop iterations take different amount of time

The execution time of a parallel region is the time of the slowest thread





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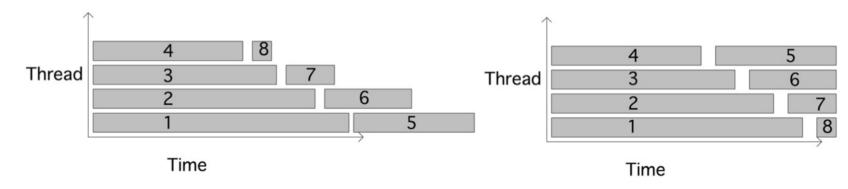
IRREGULAR WORKLOAD

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Workloads can be classified as

- Regular: each loop iteration takes the same amount of time
- Irregular: loop iterations take different amount of time

The execution time of a parallel region is the time of the slowest thread





OpenMP – Loop Scheduling

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OpenMP loop scheduling options

static

- Iterations are assigned equally among threads prior to the parallel section
- Small overhead
- Static user-defined chunk size
- dynamic
 - Each thread gets iterations according to their throughput
 - Higher overhead
 - Static user-defined chunk size

guided

- Similar to dynamic but adjusts the chunk size during runtime
- Less overhead than dynamic

	ST	ATIC SCHEDU	ILING	
Thread 0	it O	it 1		it 2
Thread 1	it 3 it	4 it	5	
Thread 2	it 6	; i	it 7	it 8
Thread 3	it 9 it	10 it 11		
	Dy	time		
Thread 0	it O	<i>it 5</i>	<i>it 10</i>	it 11
Thread 1	it 1		it 6	
Thread 2	it 2		it 8	
Thread 3	it 3 it	4 it 7	it 9	

OpenMP – Loop Scheduling

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OpenMP loop scheduling options

- static
- dynamic
- guided

Chunk size can also be user defined

- Set as a single iteration by default
- Larger chunks require less scheduling overhead
- Smaller chunks are better to schedule irregular workloads

1	#pragma omp parallel
2	{
3	<pre>#pragma omp for schedule(dynamic,10)</pre>
4	for (i=1; i <n; i++)<="" th=""></n;>
5	b[i] = (a[i] + a[i-1]) / 2.0;
6	}



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Lab Session

Vector Sum

- 1. A simple parallelization
 - Parallelize the code using **#pragma omp parallel**
 - Distribute the iterations using a **#pragma omp for** directive
 - Use the critical directive on shared data if necessary
 - Execute and measure the performance of the code
- 2. Privatize shared data
 - Privatize the accesses to shared data
 - each thread should contain it's copy of sum
 - Implement a manual merge of the partial results
 - Execute and measure the performance of the code
- 3. Removing implicit synchronizations
 - Replace the manual merge of the results with OpenMP's reduction directive
 - Execute and measure the performance of the code

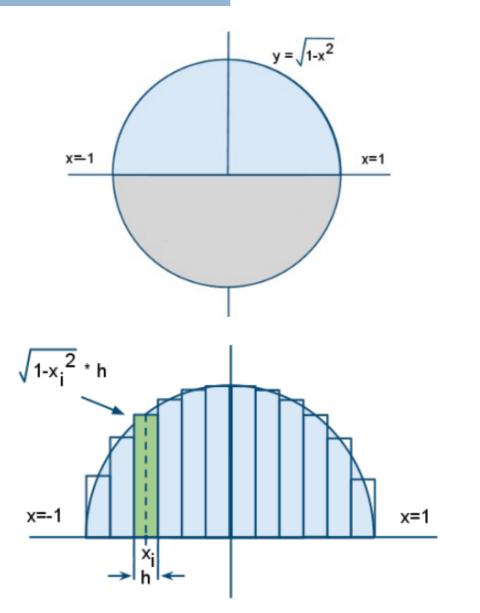
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4	<pre>int vector_sum (int array[], int size) {</pre>
5	<pre>int sum = 0;</pre>
6	
7	for (int i = 0; i < size; i++)
8	<pre>sum += array[i];</pre>
9	
0	return sum;
1	}



Pi Calculation - Integral Approximation

Pi can be calculated through the area of a circle

- The integral of the function of a circle is equivalent
- Integrals can be approximated through iterative methods



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Pi Calculation - Integral Approximation

$\sqrt{1-x_i^2}$ h x=-1 x_i x_i x_i

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For loop parallelization

- Parallelize the code using the omp parallel and omp for directives
- Privatize variables if needed
- Each thread should compute a partial sum (avoid race conditions!)
- Merge the partial results using OpenMP directives
- Execute the code and measure its performance

4	double pi_integration (long num_steps) {
5	int i;
6	double x, pi, sum = 0.0;
7	<pre>double step = 1.0 / (double) num_steps;</pre>
8	
9	for (i = 0; i < num_steps; i++) {
10	x = (i + 0.5) * step;
11	sum = sum + 4.0 / (1.0 + x * x);
12	3
13	
14	pi = step * sum;
15	
16	return pi;
17	}



Thank you for attending!

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🥑 @minhoacc